On connections between k-coloring and Euclidean k-means

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Joint Work with

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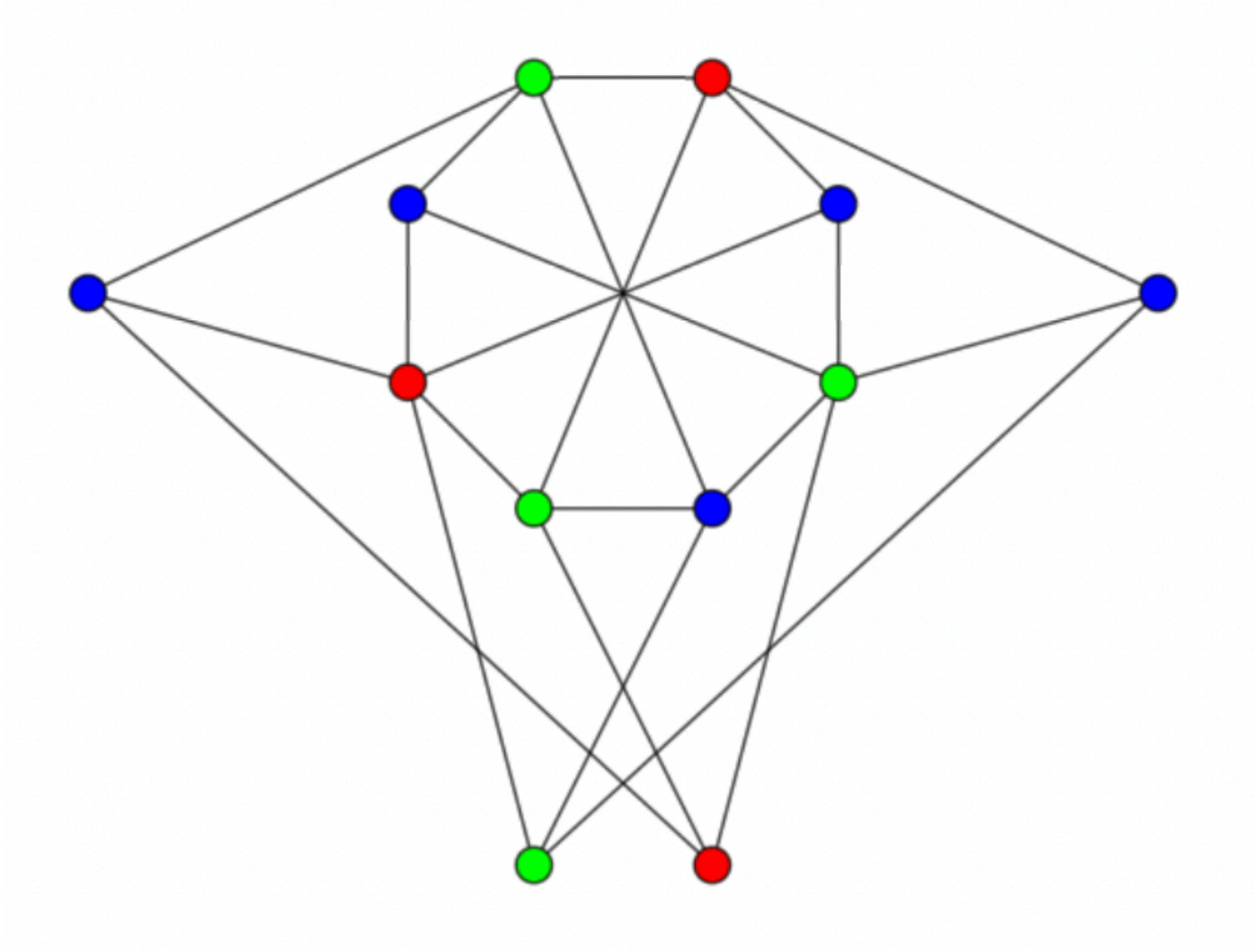
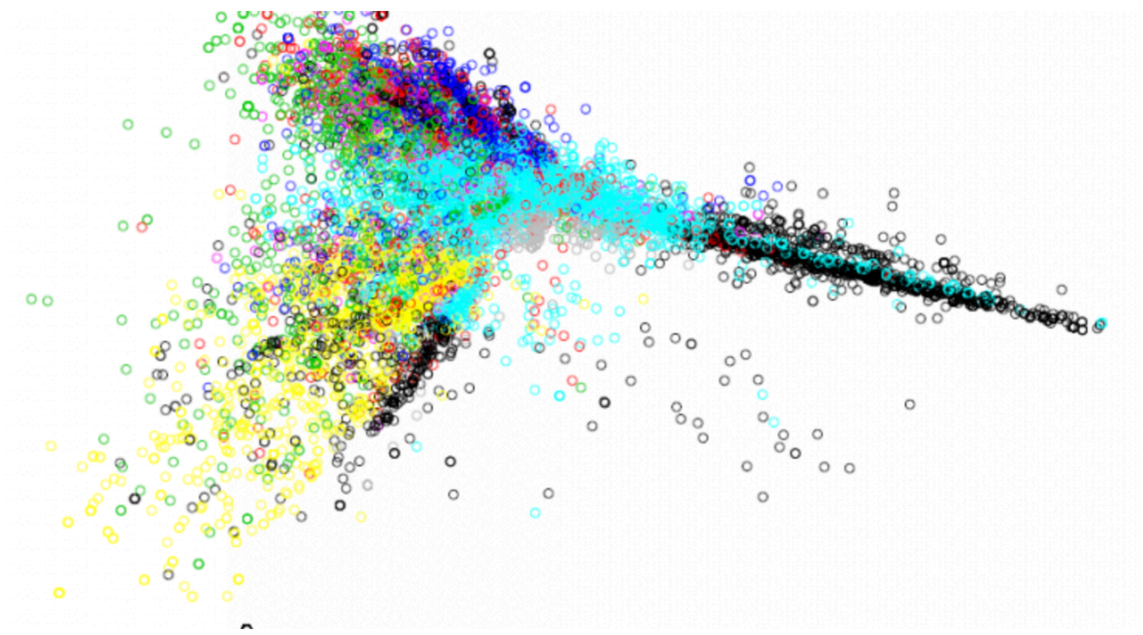
Graduated May 2023

Graph Coloring

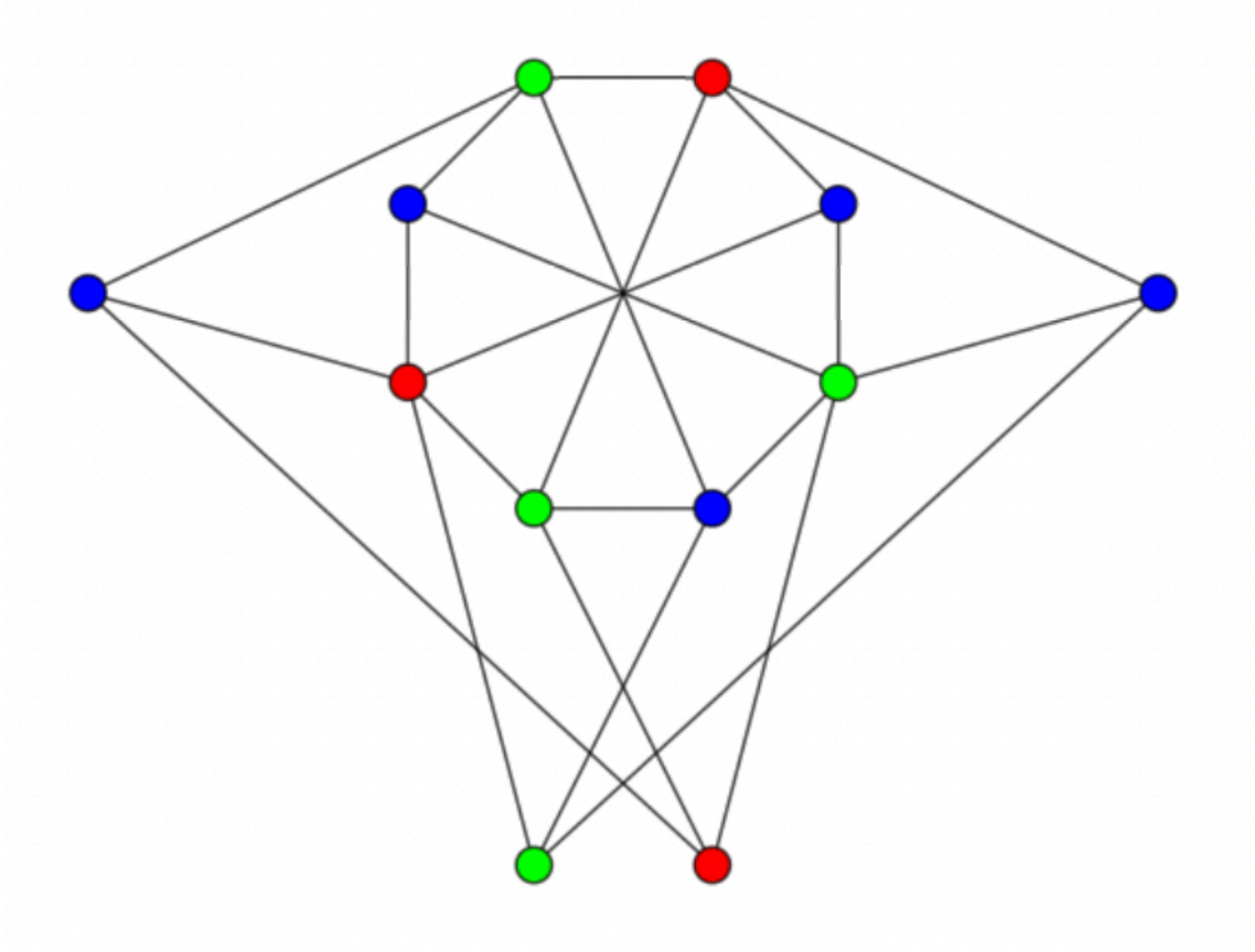


k-means Clustering



Graph Coloring

- **Input**: Graph $G = (V, E)$ and integer k



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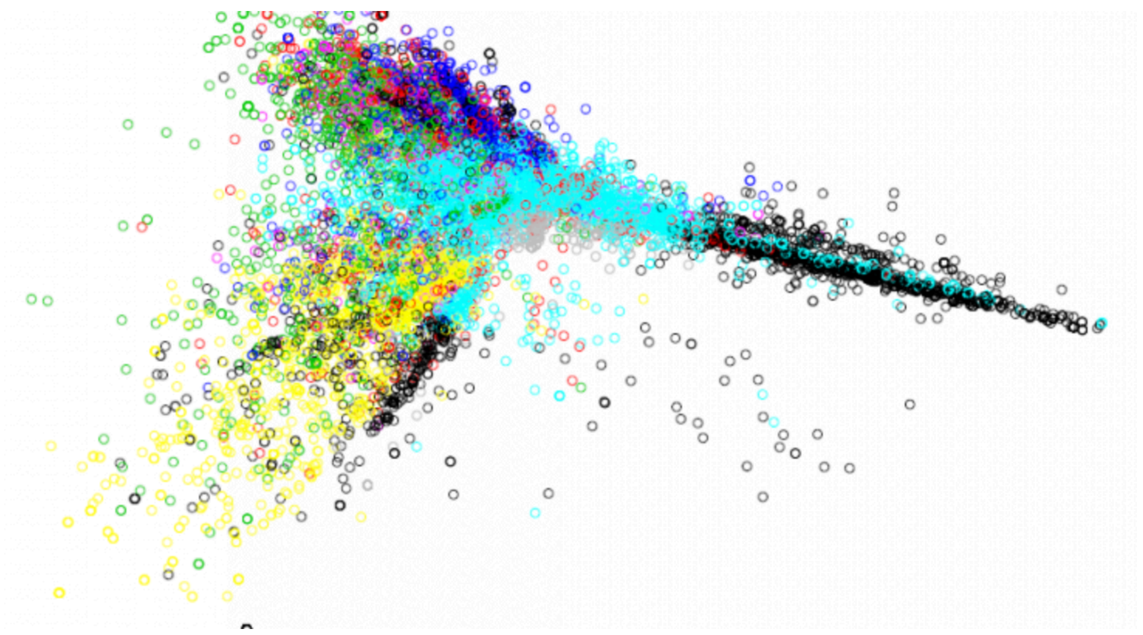
- **Input**: Points $P \subset \mathbb{R}^d$ and integer k



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- **Input**: Graph $G = (V, E)$ and integer k
- **Output**: $V := V_1 \dot\cup \cdots \dot\cup V_k$, such that

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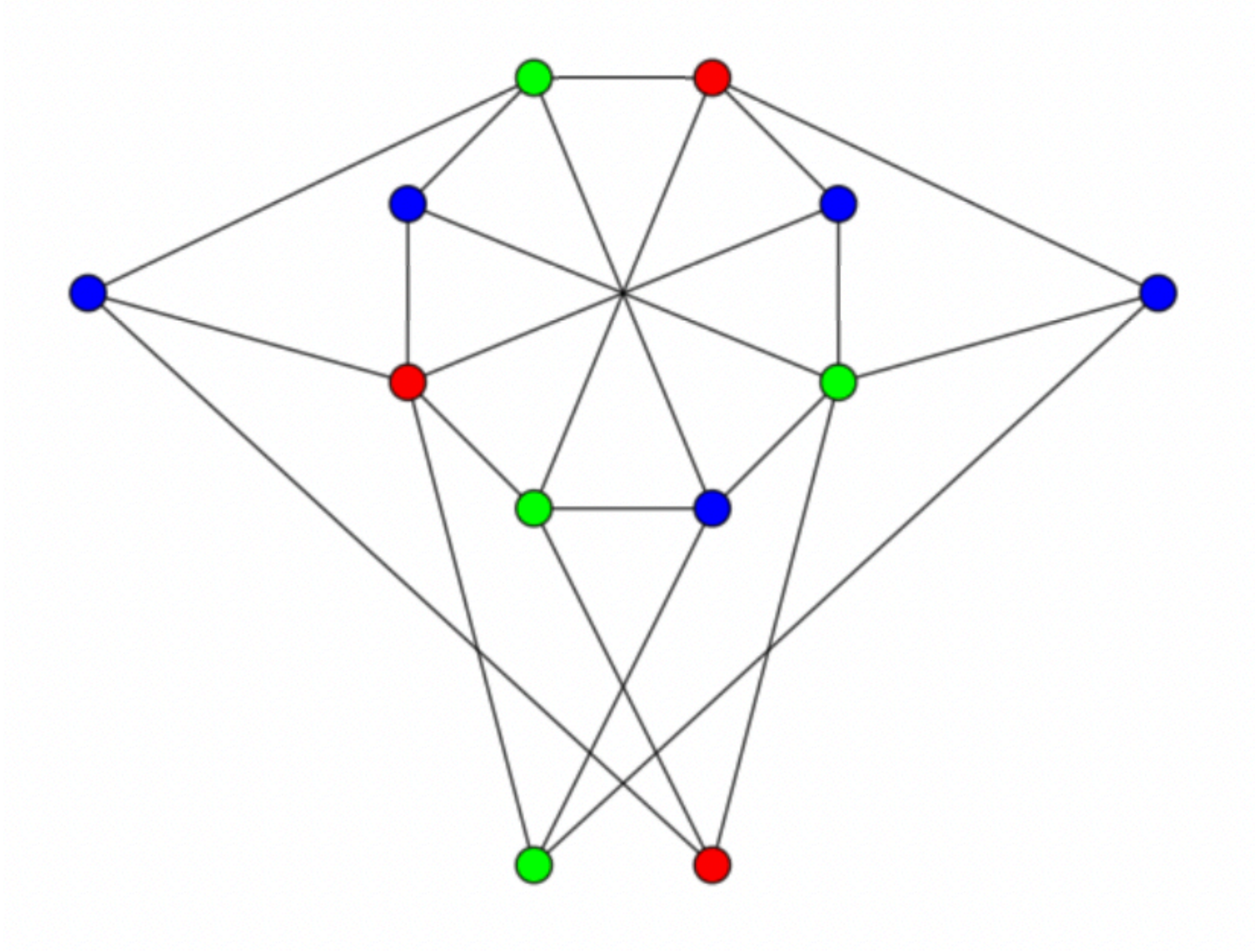

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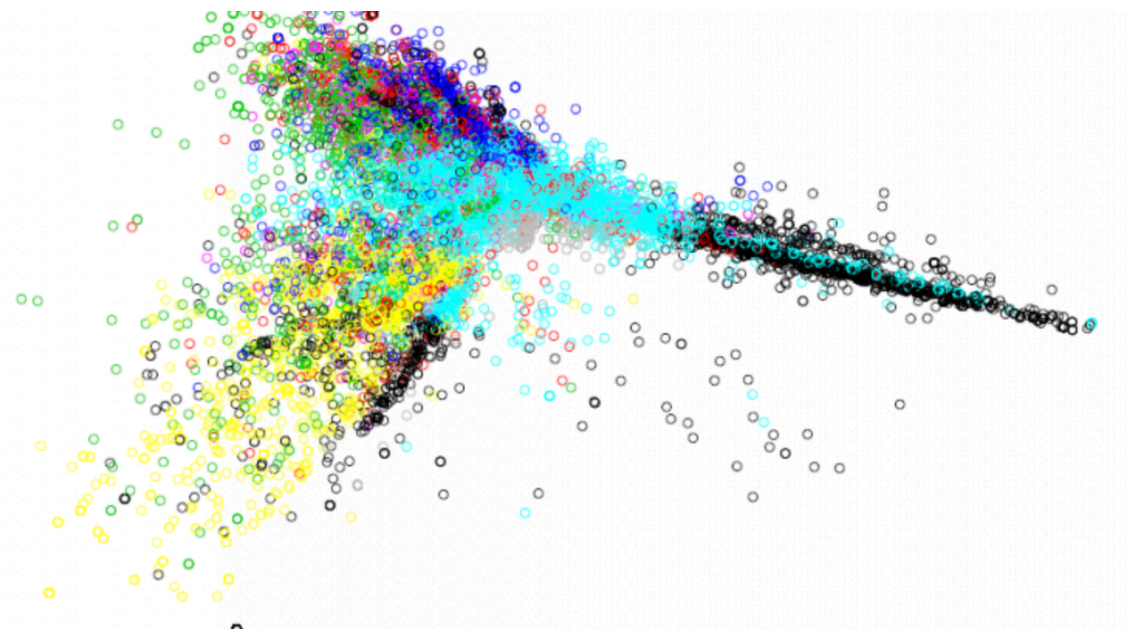
$$\sum_{i \in [k]} \sum_{p \in P_i} \|p - c_i\|_2^2 \text{ is minimized,}$$

$$\text{where } c_i = \sum_{p \in P_i} \frac{p}{|P_i|}$$



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- **Input**: Graph $G = (V, E)$ and integer k
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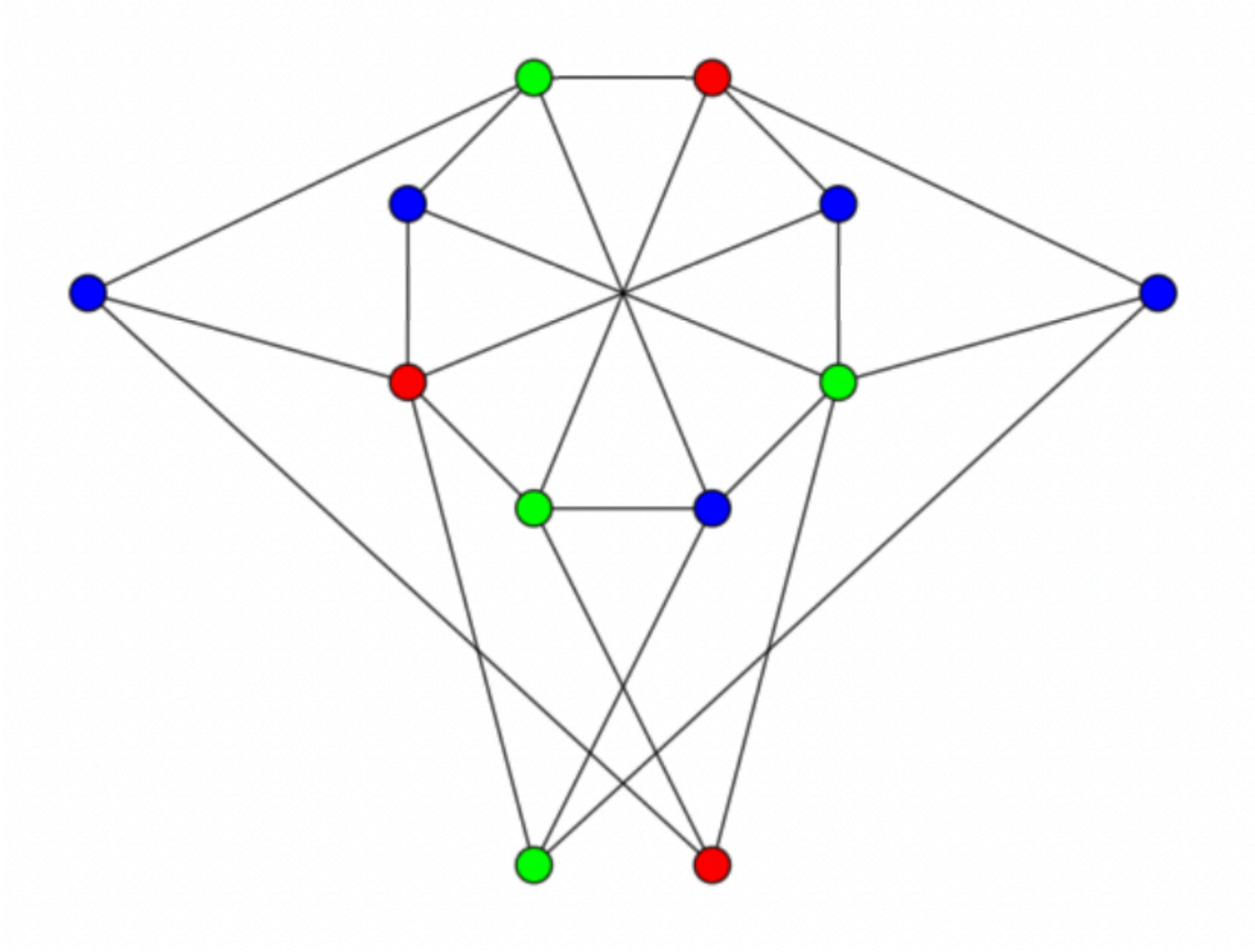
- Classic NP-hard problem



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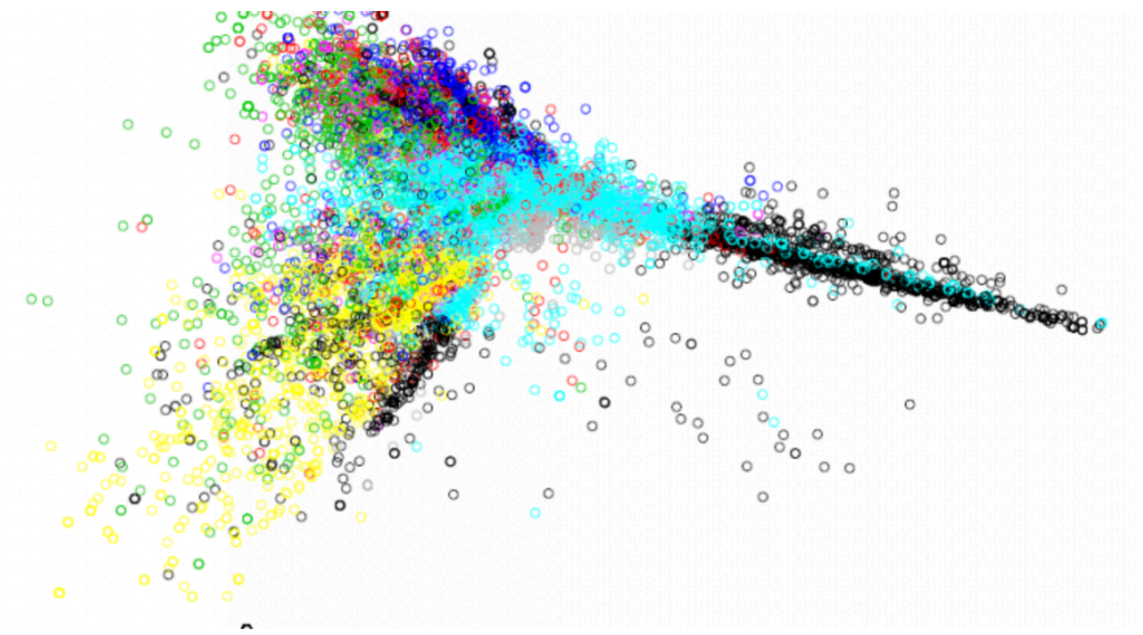
- Important NP-hard problem



k-coloring to k-means

A simple reduction

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Orient E arbitrarily

$$v \in V \longrightarrow p_v \in \mathbb{R}^{|E|}$$

$$p_v(e) = \begin{cases} +1 & \text{if } e \text{ is outgoing edge of } v \\ -1 & \text{if } e \text{ is incoming edge of } v \\ 0 & \text{otherwise} \end{cases}$$

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$\{u, v\} \in E \implies$

$\|p_v - p_u\|_2^2 = 4 + 2(d-1) = 2d + 2$

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$\{u, v\} \notin E \implies$

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Completeness

Soundness

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Completeness

- $V := V_1 \dot\cup \cdots \dot\cup V_k$, such that
  V_i is an independent set
- $\sum_{i\in[k]} \frac{1}{2|P_i|} \sum_{p,q\in P_i} \|p-q\|_2^2 = \sum_{i\in[k]} \frac{1}{2|P_i|} \sum_{p,q\in P_i} 2d$

  $= \sum_{i\in[k]} d\cdot(|P_i|-1) = d\cdot(|V|-k)$

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  $= \sum_{i\in[k]} d \cdot (|P_i| - 1) = d \cdot (|V| - k)$

Soundness

- $P := P_1 \dot\cup \cdots \dot\cup P_k$ is some clustering
- $\forall i \in [k], \sum_{p,q\in P_i} \|p-q\|_2^2 \geq 2d \cdot |P_i| \cdot (|P_i| - 1)$
- $\exists i \in [k], \sum_{p,q\in P_i} \|p-q\|_2^2 > 2d \cdot |P_i| \cdot (|P_i| - 1)$

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$$\|p_u - p_v\|_2^2 = \begin{cases} 2d & \text{if } \{u,v\} \notin E \\ 2d + 2 \;\ldots & \ldots \in E \end{cases}$$

3-coloring is NP-hard

$\Longrightarrow$

3-means is NP-hard

Completenes…

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What about 1-means?

What about 2-means?

- 2-coloring reduces to 2-means
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- **Proof Strategy**:

  (i) Structured 3-NAE-SAT is NP-hard

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  (iii) Distance matrix can be realized in Euclidean space (PSD check)

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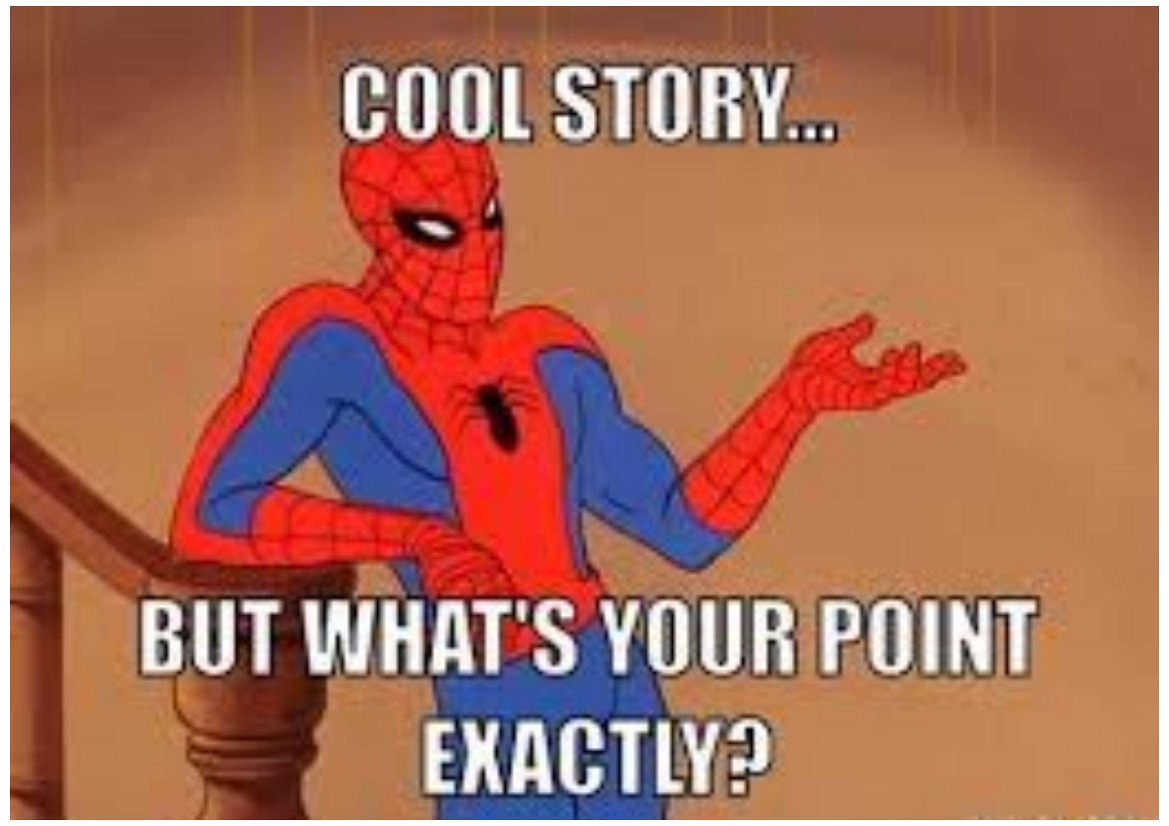

Almost 2-coloring to 2-means

New Reduction

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New Reduction

Balanced Max-Cut

- Input: Graph $G = (V, E)$, G is d-regular
- Output: $V := V_1 \dot\cup V_2$, to minimize:

$$|E \setminus E(V_1, V_2)| \cdot \left(1 + \frac{||V_1| - |V_2||}{|V|}\right)$$

2-means Clustering

- Input: Points $P \subset \mathbb{R}^d$
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Reverse the Reduction

Max-Cut on n vertices can be solved in $2^{\omega n/3}$ time

[Williams'05]

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2-means on n points can be solved in $2^{\omega n/3}$ time

This paper

Reverse the Reduction

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State-of-the-art
$2^{\omega n/3} \leq 1.73^n$

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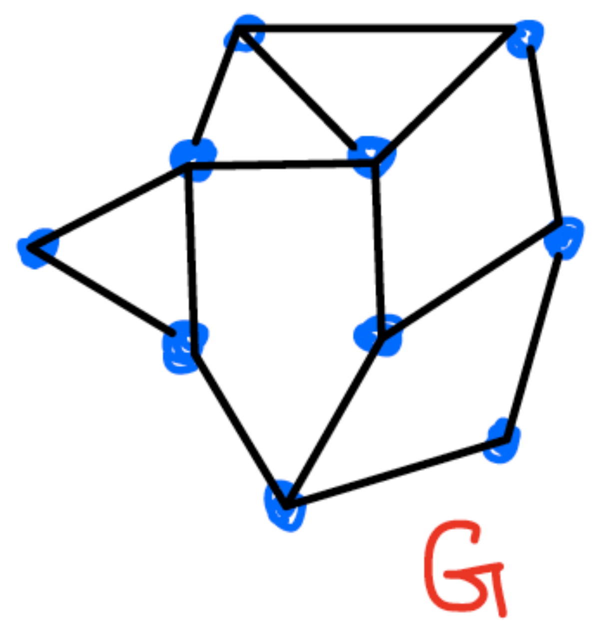

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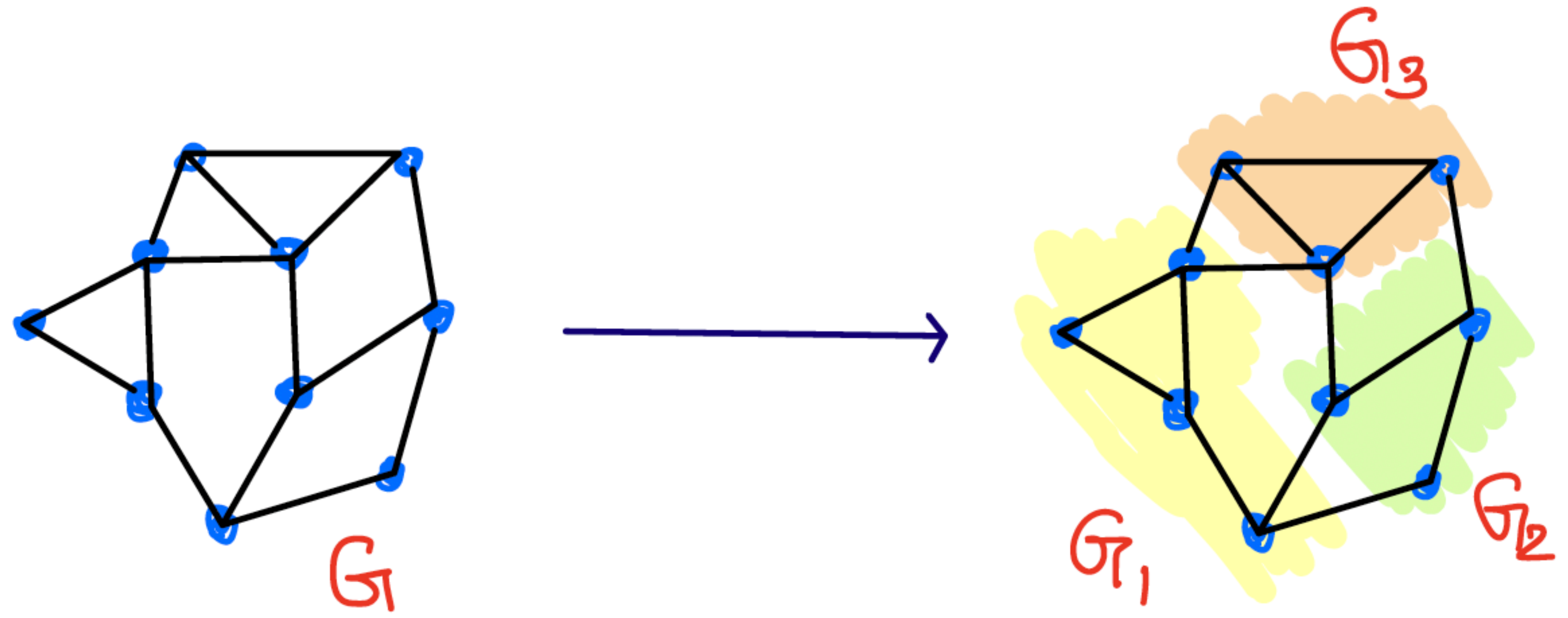

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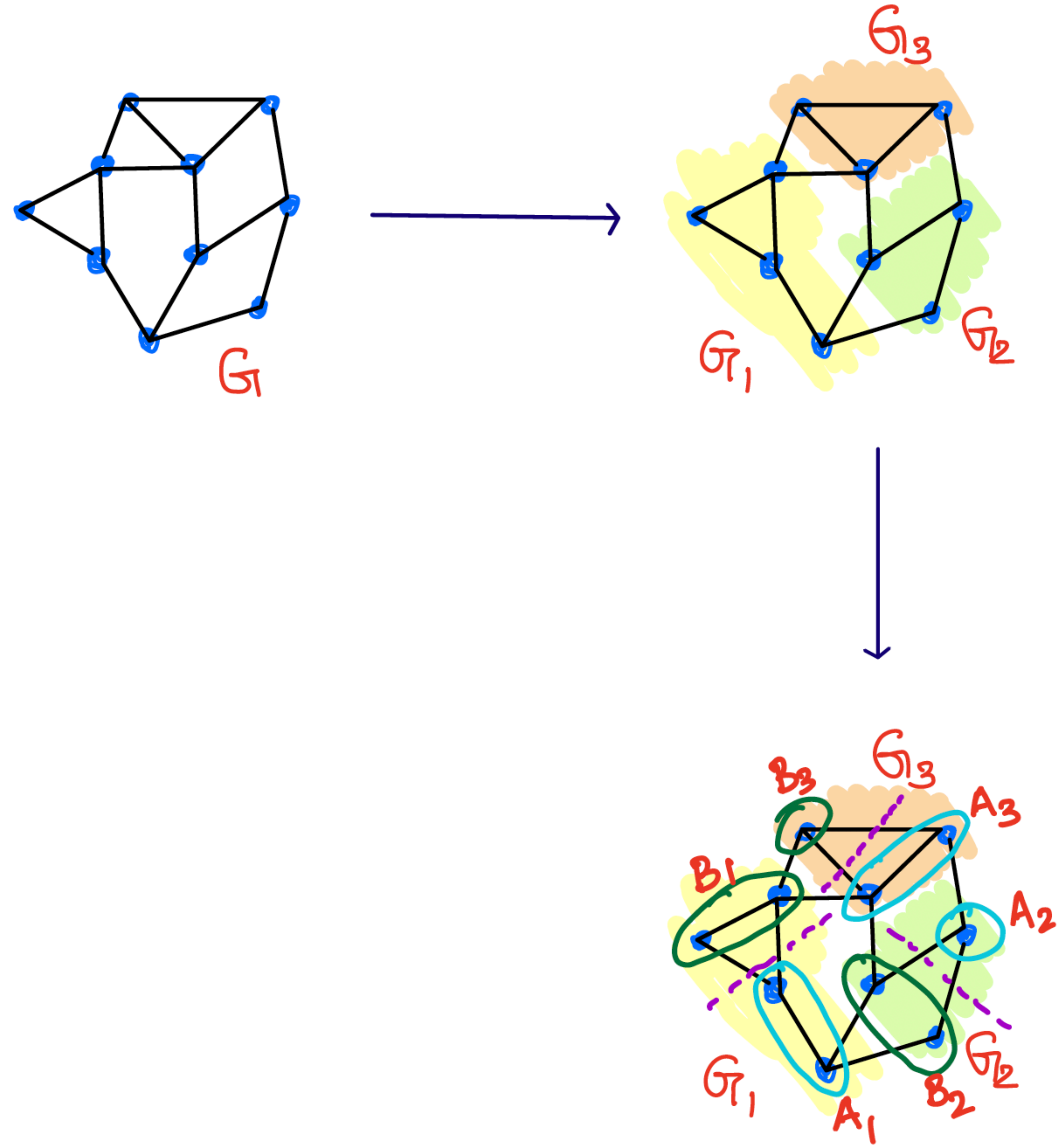
[Williams'05]

- Enumerate all cuts (A_i, B_i) of G_i



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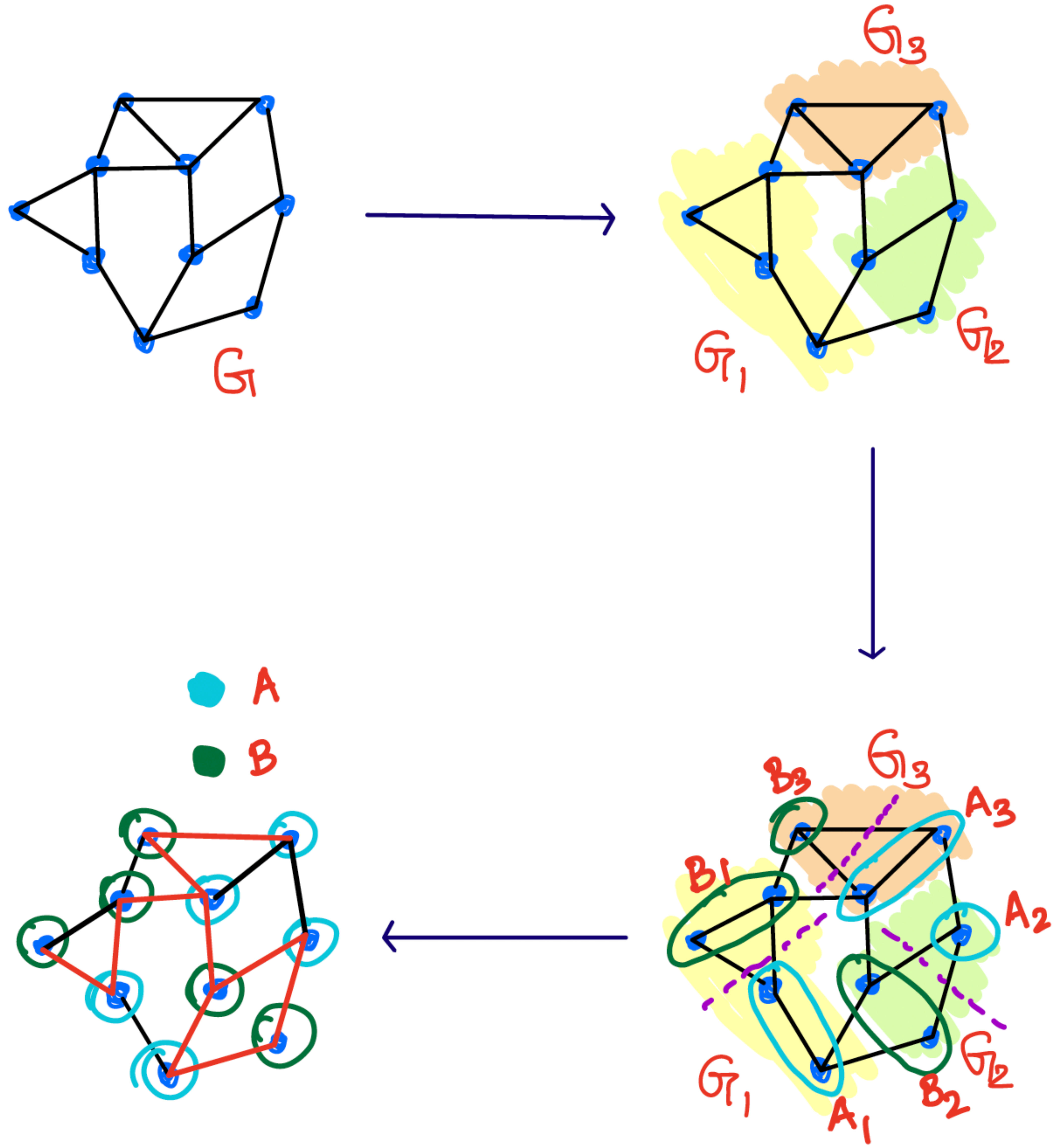
Max-Cut on n vertices can be solved in $2^{\omega n/3}$ time

[Williams'05]

- Enumerate all cuts (A_i, B_i) of G_i
- Construct graph H on $3 \cdot 2^{n/3}$ nodes
- Edge Weight of H is sum of cut edges



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- Enumerate all cuts (A_i, B_i) of G_i
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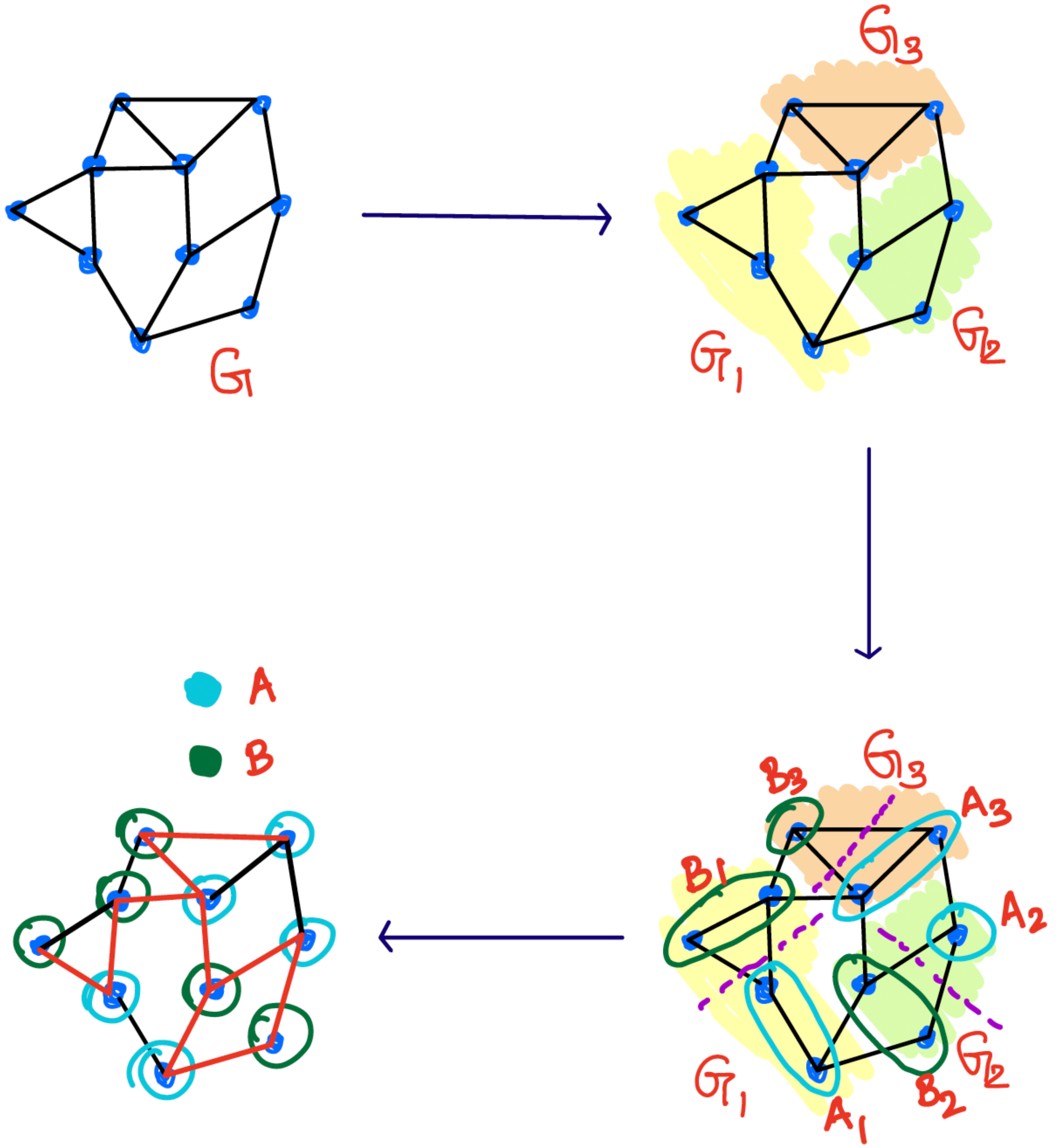
- Solve weighted triangle detection on $3 \cdot 2^{n/3}$ node graph H



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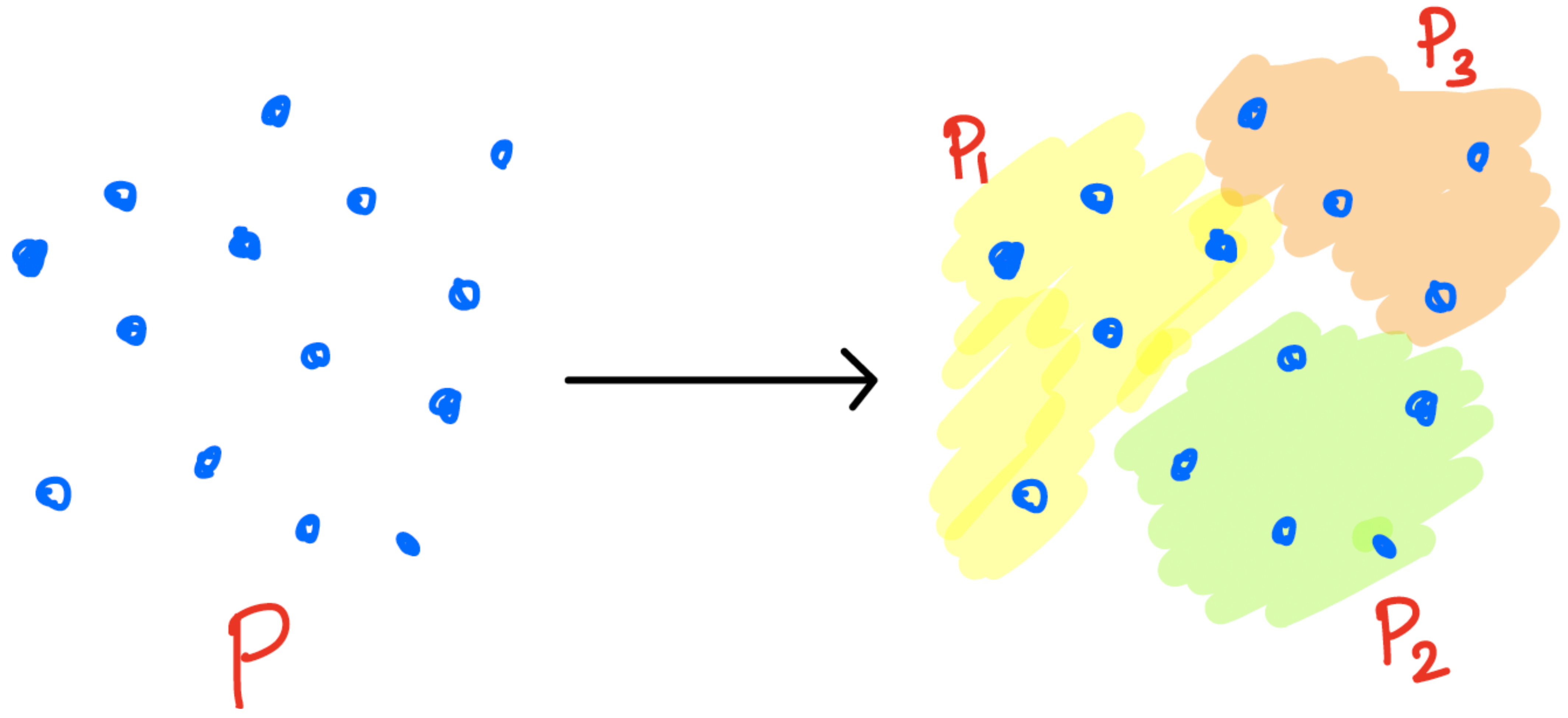

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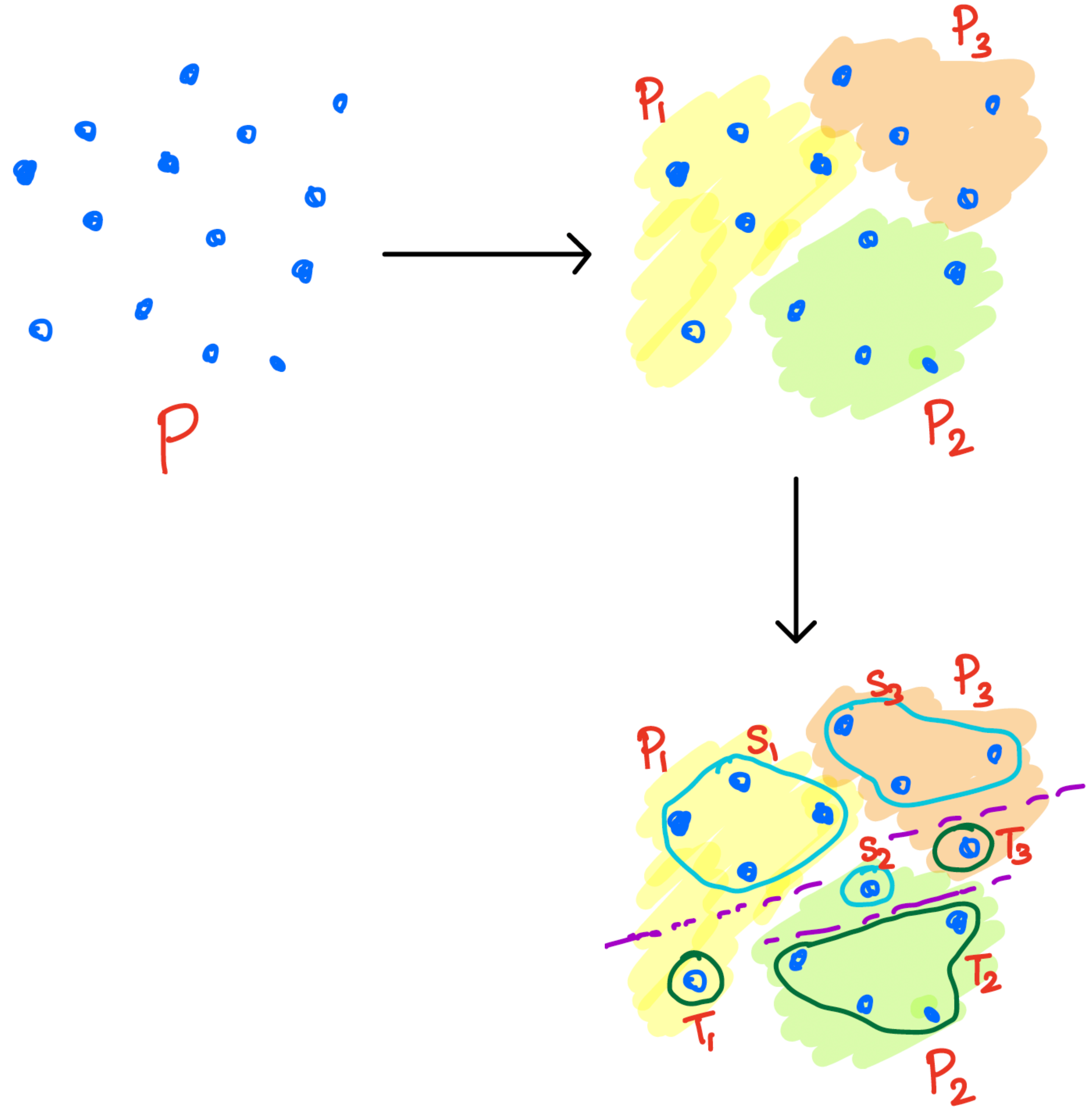

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2-means on n points can be solved in $2^{\omega n/3}$ time

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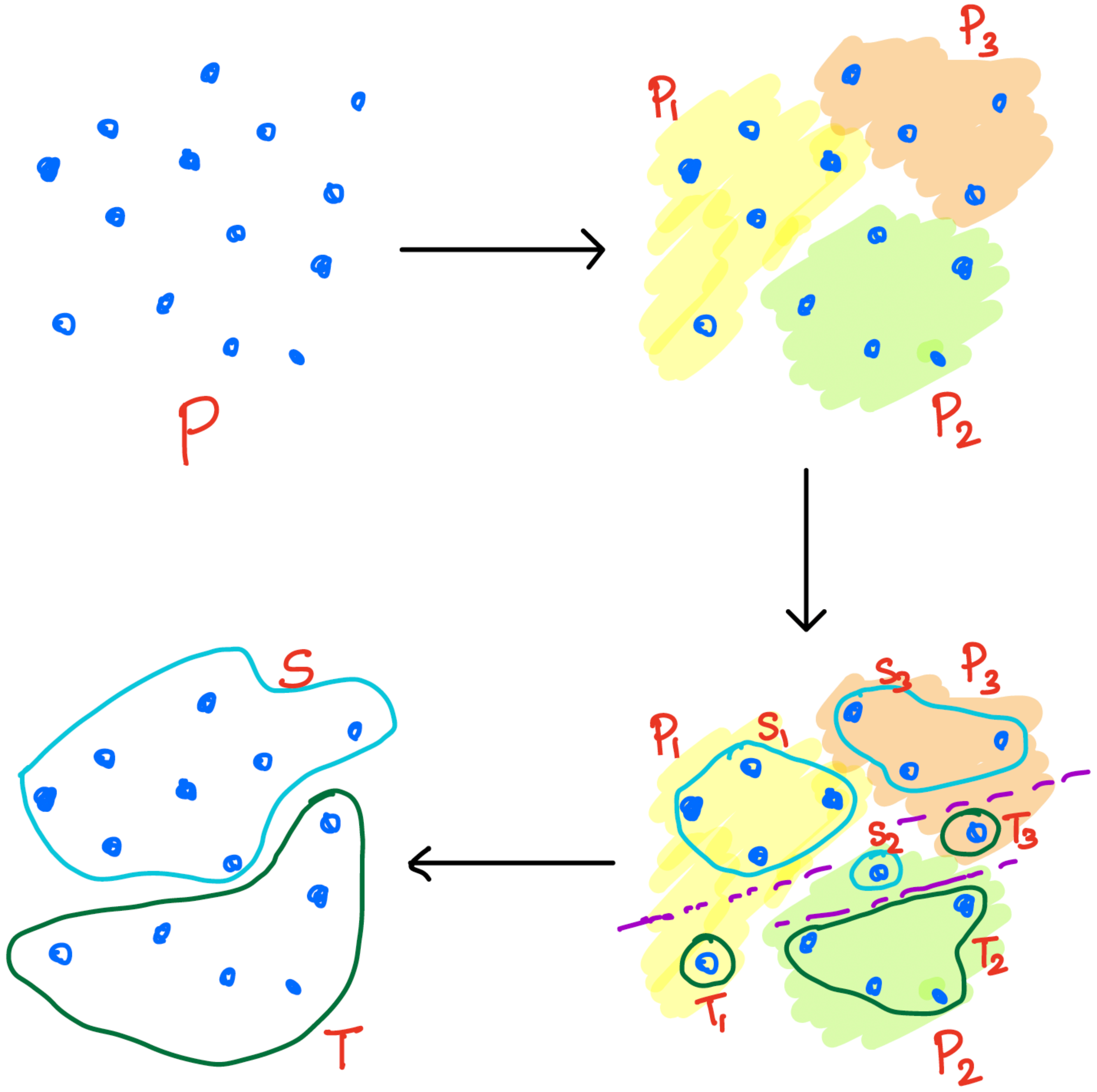

Open Directions

3-coloring $\implies$ **3**-means

Open Directions

3-coloring $\overset{?}{\Longleftarrow\!\Longrightarrow}$ **3**-means

Can we use ideas from
3-coloring algorithms
to obtain $(2 - \varepsilon)^n$ time algorithm
for 3-means?

Open Directions

Can we use ideas from
2-means algorithm
to obtain $(2 - \varepsilon)^n$ time algorithm
for 2-median or 2-center?

Thank you for engaging!